

ON BERNOULLI'S METHOD*

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Abstract. We generalize Bernoulli's classical method for finding poles of rational functions using the rational orthogonal Malmquist-Takenaka system. We show that our approach overcomes the limitations of previous methods, especially their dependence on the existence of a so-called dominant pole, while significantly simplifying the required calculations. A description of the identifiable poles is provided, as well as an iterative algorithm that can be applied to find every pole of a rational function. We discuss automatic parameter choice for the proposed algorithm and demonstrate its effectiveness through numerical examples.

Key words. Bernoulli's method, rational functions, Malmquist-Takenaka system, pole identification

MSC codes. 65T65, 30E10, 41A20

1. Introduction. Numerical methods focusing on rational approximation and interpolation have provided a rich area of research in the last decades [12, 22, 27, 37, 38, 39]. Many fields, such as control and system theory [39] and partial differential equations [12, 22, 27, 37, 38] have benefited from such approaches. In this work, we discuss the problem of finding the poles of rational functions by generalizing a method known as Bernoulli's method. As we later point out the proposed methods have great application potential especially in the field of system identification. Daniel Bernoulli considered the problem of finding the dominant (largest in absolute value) zero of a polynomial P . Identifying the zeros of the n -th degree polynomial P is equivalent to finding the poles of the rational function $R(z) := \frac{1}{z^n P(1/z)}$. Supposing that R has a unique dominant pole (the smallest in absolute value) outside the closed disk $\overline{\mathbb{D}}$, the ratios c_{n+1}/c_n constructed from the coefficients of the expansion

$$(1.1) \quad R(z) = \sum_{n=0}^{\infty} c_n z^n \quad (|z| \leq 1)$$

converge to this dominant pole [15]. In (1.1), the coefficients c_n are the Fourier-coefficients of R with respect to the trigonometric system [11].

We now proceed to give a brief historical background about this method based on the monographs [14] and [17]. Bernoulli calculated c_n in (1.1) using a recursion applied to the coefficients of P . We note that using the terminology of system theory, the Fourier-coefficients c_n can also be interpreted as the impulse response of the SISO

*Submitted to the editors on 12.10.2022.

Funding: Project no. C1748701 and K146721 have been implemented with the support provided by the Ministry of Culture and Innovation of Hungary from the National Research, Development and Innovation Fund, financed under the NVKDP-2021 and the K_23 "OTKA" funding schemes, respectively. The research was supported by the European Union within the framework of the National Laboratory for Autonomous Systems. (RRF-2.3.1-21-2022-00002).

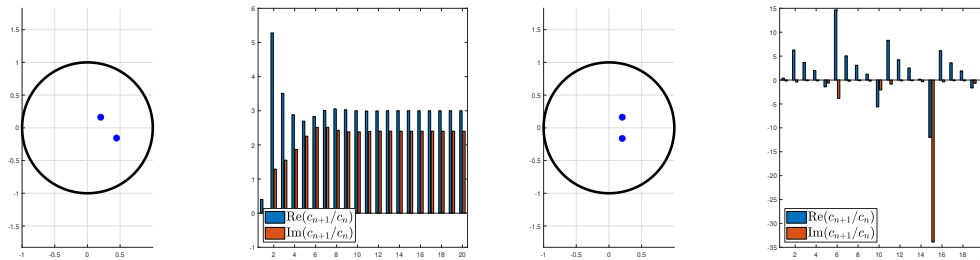
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(single input single output) [2] system whose transfer function is R . The original idea of Bernoulli was expanded by König, who generalized the pole finding method to meromorphic functions [17, 18]. Since then, many subsequent generalizations have been introduced. We note the work of Aitken, who showed that the determinants of Hankel-matrices created from the coefficients c_n can be used to approximate every pole, provided that the absolute values of the poles are pairwise different [1, 17]. Using the so-called qd (quotient-difference) algorithm, Rutishauser [29, 30] and Henrici [14] further improved Aitken's results. Detailed results on the relationship between Hankel determinants, the product of the poles and the qd algorithm can be found in chapter 7 and chapter 3 of [14] and [17], respectively.

As illustrated in Figure 1, Bernoulli's method diverges if the rational function has more than one dominant pole. We note that the above mentioned generalizations are also prone to this limitation of the method. In addition, this excludes the possibility of using Bernoulli's method for identifying the poles of SISO transfer functions, since realizable systems often have complex conjugate pairs as dominant poles.



(a) LEFT: Inverse poles (reflections of the poles across the torus) of the rational function (dominant pole uniquely exists). RIGHT: the sequence c_{n+1}/c_n .

(b) LEFT: Inverse poles of the rational function (no unique dominant pole). RIGHT: Real and imaginary parts of c_{n+1}/c_n .

Fig. 1: Bernoulli's pole finding method. The sequence c_{n+1}/c_n diverges if the function has multiple dominant poles.

In [32], a generalization of Bernoulli's method was proposed, where the discrete Laguerre-Fourier coefficients of R were considered. Using this approach we can overcome the above mentioned limitation of Bernoulli's algorithm and reconstruct a larger subset of the poles of R . In fact, the algorithm proposed in [32] can be used to reconstruct every dominant pole of R . Later, using the ideas in [32] the von Mises algorithm, which is capable of finding the dominant eigenvalues of matrices was generalized in [33]. In addition, using the so-called fartherst-point Voronoi mappings [4] induced by the pseudo-hyperbolic metric, we were able to characterize the poles of the function R which can be reconstructed by this method.

The main contribution of this work is a further generalization of the ideas proposed in [32]. Namely, we propose to use the coefficients of periodic Malmquist-Takenaka series [16, 24, 35, 39] to find the inverse poles of R . The proposed methods will include the ideas discussed in [32] as a special case. One important advantage of our generalization is that using the coefficients from a periodic Malmquist-Takenaka expansion, we can construct an iterative algorithm to find every pole of R .

The paper is organized as follows. In section 2 we discuss periodic Malmquist-

63 Takenaka systems, generalize the concept of dominant poles and introduce a general-
 64 ization of Bernoulli's algorithm. In section 3 we describe the poles which can be found
 65 using the proposed method. In section 4 we consider the problem of discretization. In
 66 section 5 we propose an iterative algorithm to identify every pole of a rational func-
 67 tion based on periodic Malmquist-Takenaka coefficients. We discuss some numerical
 68 considerations of the proposed methods in section 6, then conclude our work with an
 69 overview and future plans in section 7.

70 **2. A generalization of Bernoulli's method.** In Bernoulli's pole finding
 71 method, the coefficients c_n refer to the Fourier coefficients of the rational function R
 72 with respect to the trigonometric system $(\epsilon_n, n \in \mathbb{Z})$ [11]:

$$73 \quad c_n = \langle R, \epsilon_n \rangle := \frac{1}{2\pi} \int_0^{2\pi} R(e^{it}) e^{-int} dt \quad (n \in \mathbb{N}).$$

74 Provided that the function values of R are available on the torus, c_n can be
 75 calculated. We note that Bernoulli's method can also be applied using discrete Fourier
 76 coefficients instead of c_n . For the elementary rational functions

$$77 \quad (2.1) \quad r_\alpha(z) := \sum_{n=0}^{\infty} \bar{\alpha}^n z^n = \frac{1}{1 - \bar{\alpha}z}.$$

78 Bernoulli's algorithm can easily be verified. The number $\alpha^* := 1/\bar{\alpha}$ is the pole of the
 79 function r_α . Since α is the reflection of α^* accross the torus $\mathbb{T} := \{z \in \mathbb{C} : |z| = 1\}$,
 80 we will refer to α as the inverse pole of r_α henceforth.

81 Let \mathcal{A} denote the set of analytic functions on the closed disk. The classical
 82 Bernoulli method is summarized by the next theorem.

83 **THEOREM 2.1** (Bernoulli's algorithm). *Suppose the multiplicity of each inverse*
 84 *pole of the rational function $R \in \mathcal{A}$ is 1. If $\alpha_0 \in \mathbb{D}$ is the dominant inverse pole of R ,*
 85 *or in other words for any $\alpha \neq \alpha_0$ inverse pole, $|\alpha_0| > |\alpha|$ holds, then*

$$86 \quad (2.2) \quad \frac{\langle \epsilon_n, R \rangle}{\langle \epsilon_{n-1}, R \rangle} = \alpha_0 + O(q^n),$$

87 where $q = \max_{\alpha_0 \neq \alpha} |\alpha|/|\alpha_0|$.

88 We note that the convergence also holds when R has higher multiplicity inverse
 89 poles, however in this case the rate of convergence is $O(1/n)$.

90 The main contribution of this paper is the generalization of (2.2) to Malmquist-
 91 Takenaka systems generated by periodic sequences (or in short PMT systems), which
 92 contain the Laguerre and trigonometric systems as special cases. In this section we
 93 introduce a generalized version of (2.2), which allows us to identify a single inverse
 94 pole of a rational function. We also generalize the concept of dominant inverse poles
 95 and specify the inverse poles which can be found by the proposed method. Later
 96 in section 5 we introduce an algorithm based on the findings in this section, which
 97 will allow us to iteratively find every inverse pole of the rational function in question.
 98 Malmquist-Takenaka (or MT) systems [24, 35] can be described with the help of
 99 Blaschke factor [3]:

$$100 \quad (2.3) \quad B_a(z) := \frac{z - a}{1 - \bar{a}z} \quad (a \in \mathbb{D}, |z| \leq 1).$$

101 It is well-known [25, 31] that every Blaschke factor in

$$102 \quad (2.4) \quad \mathfrak{B} := \{\varepsilon B_a : (a, \varepsilon) \in \mathbb{D} \times \mathbb{T}\}$$

103 is a bijection on \mathbb{D} and \mathbb{T} , furthermore \mathfrak{B} forms a transformation group on \mathbb{D} with
104 respect to function composition. This group describes the congruence transformations
105 of the Bolyai-Lobachevsky geometry in the Poincaré disc model [6, 34].

106 The pseudo-hyperbolic distance

$$107 \quad (2.5) \quad \rho(a, b) := |B_a(b)| \quad (a, b \in \mathbb{D})$$

108 is a metric on \mathbb{D} which shows invariance towards Blaschke-transformations [25, 31]:

$$109 \quad \rho(T(a), T(b)) = \rho(a, b) \quad (a, b \in \mathbb{D}, T \in \mathfrak{B}).$$

110 Every sequence $\mathbf{a} = (a_n, n \in \mathbb{N}) \in \mathfrak{U} := \mathbb{D} \times \mathbb{D} \times \dots$ defines the MT-system [24, 36]
111 $\Phi^{\mathbf{a}} := \{\Phi_n^{\mathbf{a}} : n \in \mathbb{N}\}$, where

$$112 \quad (2.6) \quad \phi_n^{\mathbf{a}} := \sqrt{1 - |a_n|^2} \, r_{a_n} \prod_{j=0}^{n-1} B_{a_j} \quad (n \in \mathbb{N}, \mathbf{a} \in \mathfrak{U}).$$

113 It is well-known [16, 31] that MT-functions form a complete function system in
114 the Hardy space $H^2(\mathbb{D})$ if and only if

$$115 \quad \sum_{n=0}^{\infty} (1 - |a_n|) = \infty.$$

116 Furthermore, for any $\mathbf{a} \in \mathfrak{U}$, the function system $\Phi^{\mathbf{a}}$ is orthogonal with respect to the
117 scalar product in $H^2(\mathbb{D})$ defined as

$$118 \quad (2.7) \quad \langle f, g \rangle := \frac{1}{2\pi} \int_0^{2\pi} f(e^{it}) \overline{g(e^{it})} dt \quad (f, g \in H^2(\mathbb{D})).$$

119 We say that the MT system $\Phi^{\mathbf{a}}$ is p -periodic if there exists a number $p \in \mathbb{N}^* :=$
120 $\{n \in \mathbb{N} : n \geq 1\}$, such that $a_{n+p} = a_n$ ($n \in \mathbb{N}$). Such p -periodic sequences from \mathfrak{U} can
121 be identified with the elements of the space $\mathfrak{U}_p := \mathbb{D}^p$. Periodic MT-systems can be
122 described with p -order Blaschke-products:

$$123 \quad (2.8) \quad B_{\mathbf{a}}(z) = \prod_{j=0}^{p-1} B_{a_j}(z) \quad (\mathbf{a} \in \mathfrak{U}_p, z \in \overline{\mathbb{D}}).$$

124 Using (2.6) and (2.8) the p -periodic MT functions can be written as

$$125 \quad (2.9) \quad \phi_{kp+n}^{\mathbf{a}} = \phi_n^{\mathbf{a}} B_{\mathbf{a}}^k \quad (0 \leq n < p, k \in \mathbb{N}).$$

126 Using Cauchy's formula, we get that for any $\phi \in \mathcal{A}$ analytic function and elemen-
127 tary rational function (2.1)

128 (2.10)
$$\frac{1}{2\pi i} \int_{\mathbb{T}} \frac{\phi(z)}{z - \alpha} dz = \langle \phi, r_\alpha \rangle = \phi(\alpha) \quad (\alpha \in \mathbb{D}).$$

129 By (2.10) we can acquire simple formulas for the ratios of the PMT-Fourier co-
 130 efficients corresponding to the elementary rational function r_α . The special case
 131 of choosing a single parameter $\mathbf{a} = (a)$ ($a \in \mathbb{D}$) yields the so-called discrete La-
 132 guerre system [5, 16]. By (2.9) and (2.10) it is easy to see that in this case we have
 133 $\langle \phi_k^a, r_\alpha \rangle = \phi_0^a(\alpha) B_a^k(\alpha)$, thus

134 (2.11)
$$\frac{\langle \phi_k^a, r_\alpha \rangle}{\langle \phi_{k-1}^a, r_\alpha \rangle} = \frac{\phi_k^a(\alpha)}{\phi_{k-1}^a(\alpha)} = B_a(\alpha) \quad (k \in \mathbb{N}^*).$$

135 From (2.11), the inverse pole α can be easily acquired by the inverse B_{-a} of B_a .
 136 We now proceed to propose a pole reconstruction method similar to (2.11) in the
 137 general case, when we consider $\mathbf{a} \in \mathfrak{U}_p$ p -periodic sequences. For example by taking
 138 the indices

139 (2.12)
$$\nu_k := n + pk \quad (k \in \mathbb{N}^*, 0 < n < p, p > 1)$$

140 we get the ratios

141 (2.13)
$$\frac{\langle \phi_{\nu_k}^a, r_\alpha \rangle}{\langle \phi_{\nu_{k-1}}^a, r_\alpha \rangle} = \frac{\phi_n^a(\alpha)}{\phi_{n-1}^a(\alpha)} \quad (k \in \mathbb{N}^*, n > 0, p > 1).$$

142 In (2.13), if we choose $p = 1$ as a special case we get formula (2.11).

143 Before we can formulate our main claim, we need to generalize the concept of
 144 dominant poles. Let $A \subset \mathbb{D}$ be a finite set. We say that $\alpha_0 \in A$ is a $B_{\mathbf{a}}$ -dominant
 145 point in A , if

146 (2.14)
$$|B_{\mathbf{a}}(\alpha_0)| > |B_{\mathbf{a}}(\alpha)| \quad (\alpha \in A, \alpha \neq \alpha_0).$$

147 Using (2.13) and (2.14) we can generalize Bernoulli's method with the following
 148 theorem.

149 **THEOREM 2.2.** *Suppose that the inverse poles $\alpha \in A$ of the rational function R*
 150 *are simple and let α_0 be the $B_{\mathbf{a}}$ -dominant inverse pole of R . Then, the limit*

151 (2.15)
$$\lim_{k \rightarrow \infty} \frac{\langle \phi_{\nu_k}^a, R \rangle}{\langle \phi_{\nu_{k-1}}^a, R \rangle} = \frac{\phi_n^a(\alpha_0)}{\phi_{n-1}^a(\alpha_0)} \quad (\mathbf{a} \in \mathfrak{U}_p, p \geq 1)$$

152 *exists and the rate of convergence in (2.15) is $O(q^k)$, where*

153
$$q := \max_{\alpha \in A, \alpha \neq \alpha_0} |B_{\mathbf{a}}(\alpha)| / |B_{\mathbf{a}}(\alpha_0)|.$$

154 *Proof.* Let $R(z) := \sum_{\alpha \in A} \lambda_\alpha r_\alpha(z)$ ($z \in \mathbb{D} \cup \mathbb{T}, \lambda_\alpha \in \mathbb{C}$) be an analytic rational
 155 function on the closed disk. Then, by (2.14) $B_{\mathbf{a}}(\alpha_0) \neq 0$. Furthermore, by (2.9)
 156 and (2.10)

$$\begin{aligned}
\langle \phi_{\nu_k}^{\mathbf{a}}, R \rangle &= \bar{\lambda}_{\alpha_0} \phi_{\nu_k}^{\mathbf{a}}(\alpha_0) + \sum_{\alpha \in A \setminus \{\alpha_0\}} \bar{\lambda}_{\alpha} \phi_{\nu_k}^{\mathbf{a}}(\alpha) = \\
157 \quad &= B_{\mathbf{a}}^k(\alpha_0) \left(\bar{\lambda}_{\alpha_0} \phi_n^{\mathbf{a}}(\alpha_0) + \sum_{\alpha \in A \setminus \{\alpha_0\}} \bar{\lambda}_{\alpha} \phi_n^{\mathbf{a}}(\alpha) \frac{B_{\mathbf{a}}^k(\alpha)}{B_{\mathbf{a}}^k(\alpha_0)} \right) = \\
&= B_{\mathbf{a}}^k(\alpha_0) (\bar{\lambda}_{\alpha_0} \phi_n^{\mathbf{a}}(\alpha_0) + O(q^k)),
\end{aligned}$$

158 from which (2.15) follows directly. We note that, if $A \setminus \{\alpha_0\} = \emptyset$, then $q = 0$ and the
159 sequence (2.15) is constant. \square

160 We note that a similar statement can be formulated for inverse poles with higher
161 multiplicities, however in this case the rate of convergence cannot be guaranteed unless
162 the multiplicity of α_0 remains 1. Due to the special choice of the indices ν_k the above
163 ratios can be written as

$$164 \quad (2.16) \quad S(z) = S_n^{\mathbf{a}}(z) = \frac{\phi_n^{\mathbf{a}}(z)}{\phi_{n-1}^{\mathbf{a}}(z)} = \kappa_n \frac{z - a_{n-1}}{1 - \bar{a}_n z},$$

165 where

$$166 \quad \kappa := \kappa_n := \sqrt{(1 - |a_n|^2)/(1 - |a_{n-1}|^2)} \quad (z \in \mathbb{D}, 0 \leq n \leq p-1).$$

167 We can easily invert $w = S(z)$ with the formula

$$168 \quad (2.17) \quad z = Q(w) = Q^{\mathbf{a}}(w) = \frac{w/\kappa + a_{n-1}}{1 + \bar{a}_n w/\kappa}.$$

169 Using the limit $s(\mathbf{a}) := \lim_{k \rightarrow \infty} s_k(\mathbf{a})$, where

$$170 \quad (2.18) \quad s_k(\mathbf{a}) = s_k^R(\mathbf{a}) := \frac{\langle \phi_{\nu_k}^{\mathbf{a}}, R \rangle}{\langle \phi_{\nu_{k-1}}^{\mathbf{a}}, R \rangle} \quad (k \in \mathbb{N}^*),$$

171 we can rewrite (2.15) as

$$172 \quad (2.19) \quad \alpha_0 = Q^{\mathbf{a}}(s(\mathbf{a})).$$

173 In practice, applying formula (2.19) comes at the cost of numerical errors. The cause
174 of these errors is that in practice we can only consider the value of s_{m^*} (for some
175 finite m^* index) instead of the limit s . This error can be expressed by

$$176 \quad (2.20) \quad |Q^{\mathbf{a}}(s(\mathbf{a})) - Q^{\mathbf{a}}(s_{m^*}(\mathbf{a}))| \leq M(\mathbf{a}) \cdot |s(\mathbf{a}) - s_{m^*}(\mathbf{a})|,$$

177 where $M(\mathbf{a}) := \max_{w \in S(\mathbb{D})} |Q'(w)|$. The value of $Q'(w)$ can be expressed at a point
178 $w = S(z)$ as

$$179 \quad Q'(w) = \frac{1}{S'(z)} = \frac{(1 - \bar{a}_n z)^2}{\kappa(1 - a_{n-1} \bar{a}_n)},$$

180 from which we get

$$181 \quad (2.21) \quad M(\mathbf{a}) = \frac{(1 + |a_n|)^2}{\kappa|1 - a_{n-1}\bar{a}_n|}.$$

182 Practical ways to choose the parameter \mathbf{a} and estimate $|s(\mathbf{a}) - s_{m^*}(\mathbf{a})|$ are discussed
183 in section 6.

184 We note that in the special case $p = 1$, we choose $a_{n-1} = a_n = a$ in the above
185 formulas. We would like to highlight that choosing $\mathbf{a} = (0)$ yields $\phi_n^0(z) = z^n$, the
186 trigonometric system and in this case α_0 is the dominant inverse pole in the usual
187 sense. In this case equation (2.14) also has an obvious geometrical interpretation.

188 Considering the 1-periodic parameter sequence $\mathbf{a} = (a)$ ($a \in \mathbb{D}$) produces the
189 discrete Laguerre-system and condition (2.14) becomes

$$190 \quad (2.22) \quad \rho(a, \alpha_0) = |B_a(\alpha_0)| > |B_a(\alpha)| = \rho(a, \alpha) \quad (\alpha \in A),$$

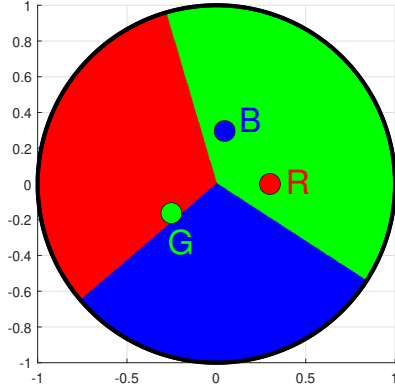
191 where $\rho(\cdot, \cdot)$ is the pseudo-hyperbolic metric as given in (2.5). We discuss the geometric
192 interpretation of (2.22) in [32]. We note that in this special case $n = 1$, $a_0 = a_1 = a$,
193 therefore

$$194 \quad M(a) = \frac{(1 + |a|)^2}{1 - |a|^2} = \frac{1 + |a|}{1 - |a|}.$$

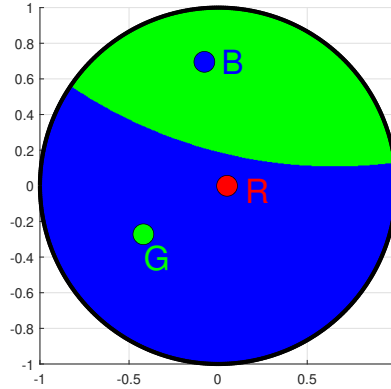
195 **3. Geometric properties of dominant poles.** In this section we summarize
196 the geometric interpretations of the generalized dominant poles (2.14). From the
197 point of view of our proposed pole identification scheme built around theorem 2.2,
198 the results in this section help us visualize how to choose the parameters of the
199 Malmquist-Takenaka expansions to identify specific poles of the rational function R .
200 Formally, using the concept of Voronoi-mappings [4], we describe some regions of \mathbb{D} .
201 Choosing the parameters of the aforementioned MT-systems from these regions and
202 applying theorem 2.2 will allow for finding specific poles of R . The results discussed
203 here also show, that some poles may be "hidden" in the sense, that independent of our
204 choice of the parameter vector \mathbf{a} , they will never be dominant (will not satisfy (2.14)),
205 therefore cannot be recovered directly using theorem 2.2. In order to find such hidden
206 poles with the proposed method, "cancelling the effect" of other poles is necessary.
207 We discuss such techniques in section 5.

208 We begin by considering the $p = 1$ case, that is, when the periodic Malmquist-
209 Takenaka system in theorem 2.2 depends on a single $a \in \mathbb{D}$ parameter. We are going
210 to illustrate that in this case, the dominant poles of R can be described using the
211 so-called pseudo hyperbolic metric and Voronoi mappings. Moreover, we are going to
212 investigate an interesting relationship between these dominant, or "visible" poles (the
213 ones that we can recover using theorem 2.2), and the extreme points of the convex
214 hull of the inverse poles. This observation will allow us to point out some interesting
215 relationships between Voronoi-mappings generated by different types of metrics and
216 corresponding variants of convex hulls. The notion of dominant poles as introduced
217 in [32] can be geometrically described using farthest-point Voronoi-mappings [4]. Let
218 F_A denote the union of the hyperbolic bisectors $\ell_{a,b} := \{z \in \mathbb{D} : \rho(a, z) = \rho(b, z)\}$ and
219 let $D_A := \mathbb{D} \setminus F_A$ ($a, b \in A, a \neq b$). Then, for each $a \in D_A$ there uniquely exists a
220 point $\alpha \in A$ which is farthest from a in metric ρ . Let $V = V_A := D_A \rightarrow A$ denote the
221 function which maps every $a \in D_A$ to the point in A farthest away from it. Then, V_A

222 is the farthest-point Voronoi mapping of the set A generated by the pseudo hyperbolic
 223 distance. The $V_A^{-1}(\alpha)$ ($\alpha \in A$) Voronoi-cells provide a disjoint partitioning of the set
 224 $\mathbb{D} \setminus A$. Condition (2.22) is equivalent to $V_A(a) = \alpha_0$, that is the set of inverse poles
 225 which are B_a ($a \in \mathbb{D}$) dominant is exactly the range of the Voronoi mapping V_A . Any
 226 inverse poles α for which $V_A^{-1}(\alpha) = \emptyset$ cannot be retrieved.



(a) Farthest-point Voronoi diagram using pseudo hyperbolic distance. In this case there were more than one dominant inverse poles in the classical sense, however using the proposed approach (2.15), we can reconstruct any of them.



(b) Farthest-point Voronoi diagram using pseudo hyperbolic distance. Not every $a \in A$ is guaranteed to have a nonempty Voronoi cell. In this case the point labeled "R" has no corresponding region.

Fig. 2: Some example Voronoi cells generated by the pseudo hyperbolic distance. Members of the set A are denoted by points of different colors and are labeled with the letters "R", "G" and "B" for red, green and blue. The corresponding Voronoi cells are shown in the same color. If A contains the inverse poles of a rational function R , then choosing the parameter of the 1-periodic MT system (discrete Laguerre system) from the set $V_A^{-1}(\alpha)$ ($\alpha \in A$) allows for the reconstruction of the inverse pole α with (2.15).

227 In figure 2, we illustrate some farthest-point cells $V_A^{-1}(\alpha)$. For the examples
 228 in figure 2, not considering points strictly on the border between two neighbouring
 229 Voronoi cells, the limit (2.15) exists for any $a \in \mathbb{D}$ parameter. The rate of convergence
 230 depends on the choice of the parameter a chosen from the Voronoi-cells. The choice
 231 of this parameter will be further discussed in section 6. Suppose that A contains the
 232 inverse poles of a rational function. Then, the examples in figure 2 also illustrate that
 233 if there is no dominant inverse pole (the points in A fall on a circle), each inverse pole
 234 can still be found using the proposed algorithm.

235 We note that in the Euclidean plane (when we define the distance generating the
 236 Voronoi mappings as $\rho(a, b) = \|a - b\|_2$ ($a, b \in \mathbb{R}^2$) instead of (2.5)), we can describe
 237 V_A using convex geometry. Namely, the range of V_A can be described by the set of
 238 extreme points of the convex hull of A (see figure 4). The analogous statement does
 239 not hold for the hyperbolic case. In figure 4, we illustrate that the set of vertices
 240 of the hyperbolic convex hull of A is larger than the range of V_A . In this case, one

241 can describe the range of V_A on the hyperbolic plane using the notion of paracyclic
 242 convexity. We plan to investigate this phenomena in detail in a future work.

243 Next, we would like to extend the idea of describing the dominant (or "visible")
 244 inverse poles of the rational function R for identification by periodic MT-systems,
 245 where $p \geq 2$. In this case, the MT-Fourier coefficients used to identify the dominant
 246 inverse poles depend on a p dimensional parameter vector denoted by \mathbf{a}_p . We are
 247 interested in describing the Voronoi cells generated by the inverse poles of R , where
 248 instead of the hyperbolic metric discussed above, the notion of distance is given by
 249 the Blaschke product corresponding to \mathbf{a}_p . The case, when the first $p - 1$ components
 250 in \mathbf{a}_p are chosen as the inverse poles $\alpha_0, \dots, \alpha_{p-2} \in A$ will be of special interest to
 251 us (see figure 3 and section 5), however we discuss our findings for a general choice
 252 of \mathbf{a} . In order to give a geometric description of the general case, let us consider the
 253 sequence $\mathbf{a} := (a_0, a_1, \dots) \in \mathfrak{U}$ and fixing the first $p - 1$ components define

$$254 \quad (3.1) \quad \mathbf{a}_p := (a_0, a_1, \dots, a_{p-2}, a) \quad (a \in \mathbb{D}).$$

255 We are going to use the vector \mathbf{a}_p to construct a periodic MT system. Then, as stated
 256 in (2.14), we call $\alpha_0 \in A$ the $B_{\mathbf{a}_p}$ -dominant element in A if for the mapping

$$257 \quad (3.2) \quad \rho_p(a, \alpha) := |B_{\mathbf{a}_p}(\alpha)| \quad (\alpha \in \mathbb{D}),$$

258 the statement analogous to (2.14) holds:

$$259 \quad (3.3) \quad \rho_p(a, \alpha_0) > \rho_p(a, \alpha) \quad (\alpha \in A, \alpha \neq \alpha_0).$$

260 Using the mappings ρ_p we can introduce the Voronoi mappings $V_{A,p}$ generated
 261 by them. For any interior point of the Voronoi cells, the limit (2.15) exists. We note
 262 that if $\alpha \in A$ and α is also a component in \mathbf{a}_p , then $V_{A,p}^{-1}(\alpha) = \emptyset$, or in other words α
 263 cannot be found using the proposed method. As discussed in detail in section 5, this
 264 property can be exploited to construct an iterative algorithm which finds every pole
 265 of the rational function. An example mapping $V_{A,p}$ is provided for $p = 2$ in figure 3.

266 **4. Discrete Malmquist-Takenaka systems.** In numerical calculations in-
 267 stead of the Fourier coefficients $\widehat{f}(n)$ [11] of a function $f : \mathbb{T} \rightarrow \mathbb{C}$ we often consider
 268 the N -periodic discrete Fourier-coefficients

$$269 \quad (4.1) \quad \widehat{f}_N(n) := \frac{1}{N} \sum_{z \in \mathbb{T}_N} f(z) z^{-n} \quad (n, N \in \mathbb{N} \setminus \{0\}),$$

270 where

$$271 \quad \mathbb{T}_N := \{e^{2i\pi k/N} : 0 \leq k < N\}.$$

272 The (finite) trigonometric system is orthogonal with respect to the discrete inner
 273 product [11]

$$274 \quad [f, g]_N := \frac{1}{N} \sum_{z \in \mathbb{T}_N} f(z) \overline{g(z)} \quad (N = 2, 3, \dots).$$

275 Let $\widehat{r}_\alpha(n)$ denote the n -th Fourier coefficient of the elementary rational function
 276 r_α . From the formula

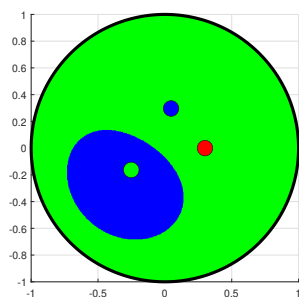
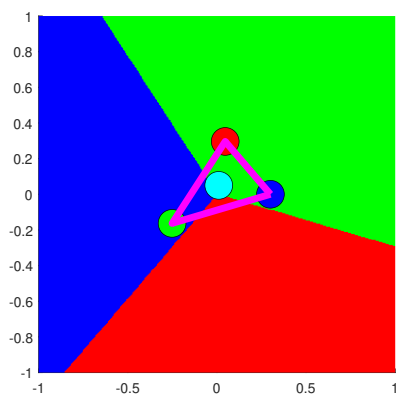
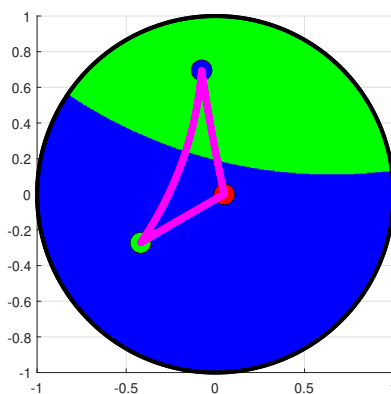


Fig. 3: An example of $V_{A,p}$ for $p = 2$. For $p > 1$, the borders between the Voronoi cells can no longer be described with hyperbolic lines. In this case a_0 is chosen as the inverse pole denoted by the red point (hence its corresponding Voronoi cell is empty). If we choose the parameter a_1 from either of the two Voronoi cells and apply theorem 2.2, then we can find the inverse pole corresponding to the color of the chosen cell.



(a) Farthest-point Voronoi mapping on the Euclidean plane (generated by Euclidean distance). The set of vertices of the convex hull of A is the range of V_A .



(b) Farthest-point Voronoi mapping generated by the pseudo hyperbolic distance. Vertices of the hyperbolic convex hull form a larger set than the range of V_A .

Fig. 4: Relationship between convex geometry and nonempty farthest-point Voronoi cells in the Euclidean and hyperbolic cases.

$$r_\alpha(z) = \sum_{n=0}^{\infty} \bar{\alpha}^n z^n \quad (\alpha \in \mathbb{D}, |z| < 1)$$

278 it follows that the discrete Fourier coefficients of the elementary rational function r_α
 279 can be written as

$$(4.2) \quad \widehat{r}_{\alpha, N}(n) = \frac{\overline{\alpha}^n}{1 - \overline{\alpha}^N} = \frac{\widehat{r}_{\alpha}(n)}{1 - \overline{\alpha}^N} \quad (0 \leq n < N).$$

Because of (4.2), we can use discrete Fourier coefficients to construct the ratios in (2.2).

In order to formulate the discrete Malmquist-Takenaka system, let us consider an N -periodic MT system $\phi_{n+Nk}^{\mathbf{a}} = \phi_n^{\mathbf{a}} B_{\mathbf{a}}^k$ generated by the vector $\mathbf{a} \in \mathbb{D}^N$. Taking the MTF coefficients of r_{α} and considering (2.10) leads to

$$(4.3) \quad \langle \phi_{n+Nk}^{\mathbf{a}}, r_{\alpha} \rangle = \phi_n^{\mathbf{a}}(\alpha) B_{\mathbf{a}}^k(\alpha) \quad (k \in \mathbb{N}, 0 \leq n < N).$$

Because $|B_{\mathbf{a}}(\alpha)| < 1$, the Malmquist-Takenaka Fourier series with the coefficients (4.3) is absolutely and uniformly convergent on \mathbb{T} . Furthermore (since the MT system is complete in the Hardy space $H^2(\mathbb{D})$) the series produces r_{α} :

$$(4.4) \quad r_{\alpha}(z) = \sum_{n=0}^{N-1} \sum_{k=0}^{\infty} \overline{\phi}_n^{\mathbf{a}}(\alpha) \phi_n^{\mathbf{a}}(z) \overline{B}_{\mathbf{a}}^k(\alpha) B_{\mathbf{a}}^k(z) = \frac{1}{1 - \overline{B}_{\mathbf{a}}(\alpha) B_{\mathbf{a}}(z)} \sum_{n=0}^{N-1} \overline{\phi}_n^{\mathbf{a}}(\alpha) \phi_n^{\mathbf{a}}(z).$$

Taking the limit $\alpha \rightarrow w \in \mathbb{T}$ in (4.4) produces the Christoffel-Darboux formula for Malmquist-Takenaka systems (see also [5, 7, 28]):

$$(4.5) \quad \sum_{n=0}^{N-1} \overline{\phi}_n^{\mathbf{a}}(w) \phi_n^{\mathbf{a}}(z) = \frac{1 - \overline{B}_{\mathbf{a}}(w) B_{\mathbf{a}}(z)}{1 - \overline{w}z} \quad (w, z \in \overline{\mathbb{D}}, w \neq z).$$

In order to acquire the discrete MT functions let us consider the set

$$(4.6) \quad \mathbb{T}_N^{\mathbf{a}} := \{z \in \mathbb{T} : B_{\mathbf{a}}(z) = 1\},$$

where $\mathbf{a} \in \mathbb{D}^N$. Since $B_{\mathbf{a}} : \mathbb{T} \rightarrow \mathbb{T}$ is an N -fold mapping [25, 31], the number of elements in $\mathbb{T}_N^{\mathbf{a}}$ is exactly N . We note that (4.6) can also be an appropriate choice of discretization points for a periodic MT system, whose period is less than N . For example, if we consider the 1-periodic (discrete Laguerre) system, choosing the discretization points (4.6) with $\mathbf{a} = (a, a, a, \dots) \in \mathbb{D}^N$ ($N \geq 1$) is appropriate. Furthermore,

$$(4.7) \quad \sum_{n=0}^{N-1} \overline{\phi}_n^{\mathbf{a}}(w) \phi_n^{\mathbf{a}}(z) = \begin{cases} 0 & (z, w \in \mathbb{T}_N^{\mathbf{a}}, z \neq w) \\ \sigma^2(z) & (z \in \mathbb{T}_N^{\mathbf{a}}, z = w), \end{cases}$$

where

$$(4.8) \quad \sigma(z) := \sum_{n=0}^{N-1} \frac{1 - |a_n|^2}{|1 - \overline{a}_n z|^2}.$$

From the equations (4.7) and (4.8) it follows that the function system $\phi_n^{\mathbf{a}}$ ($0 \leq n < N$) is orthonormal with respect to the inner product

$$307 \quad (4.9) \quad [f, g]_N^{\mathbf{a}} := \sum_{z \in \mathbb{T}_N^{\mathbf{a}}} f(z) \bar{g}(z) / \sigma(z),$$

308 or in other words $[\phi_n^{\mathbf{a}}, \phi_m^{\mathbf{a}}]_N^{\mathbf{a}} = \delta_{mn}$. Since $\phi_{n+kN}^{\mathbf{a}} = \phi_n^{\mathbf{a}}$ holds in any $z \in \mathbb{T}_N^{\mathbf{a}}$ point,
 309 the discrete MTF coefficients are N -periodic. Using this and (4.4) we can arrive at
 310 a formula analogous to (4.2) for MTF coefficients:

$$311 \quad (4.10) \quad [\phi_n^{\mathbf{a}}, r_{\alpha}]_N^{\mathbf{a}} = \frac{\bar{\phi}_n^{\mathbf{a}}(\alpha)}{1 - \bar{B}_{\mathbf{a}}(\alpha)} = \frac{\langle \phi_n^{\mathbf{a}}, r_{\alpha} \rangle}{1 - \bar{B}_{\mathbf{a}}(\alpha)} \quad (0 \leq n < N).$$

312 By (4.10) we can also use discrete MTF coefficients to construct the ratios needed
 313 for the proposed pole finding method (2.15).

314 **5. Finding every pole of a rational function.** In this section we are going
 315 to propose an iterative algorithm based on theorem 2.2, which allows for finding every
 316 inverse pole of a rational function R . As before, we are going to assume that every
 317 inverse pole is simple and denote the (finite) set of inverse poles by $A \subset \mathbb{D}$.

318 We begin by introducing a mechanism to eliminate inverse poles which have al-
 319 ready been found. From section 3 it is clear that using p -periodic MTF coefficients
 320 in (2.15), where the MT system is generated by $\mathbf{a} \in \mathbb{D}^p$ allows for the identification
 321 of a single $B_{\mathbf{a}}$ -dominant inverse pole. Modifying the parameter vector \mathbf{a} lets us find
 322 different inverse poles from A , however not every inverse pole can be acquired in this
 323 way (see figure 3). We will make use of the following observation:

$$324 \quad (5.1) \quad B_{\mathbf{a}}(\alpha) = 0 \quad (\mathbf{a} := (a_0, \dots, a_{p-2}, \alpha) \in \mathbb{D}^p, p \geq 1).$$

325 In effect (5.1) states that if the inverse pole $\alpha \in A$ is also a component of \mathbf{a} , then α
 326 cannot be $B_{\mathbf{a}}$ -dominant. This provides an opportunity to "eliminate" already found
 327 inverse poles. Suppose we applied theorem 2.2 with a PMT system defined by $\mathbf{a} \in \mathbb{D}^p$
 328 to identify the inverse pole $\alpha \in \mathbb{D}$. Now applying (2.15) using the MTF coefficients
 329 determined by the vector $\mathbf{b} := (\mathbf{a}, \alpha) \in \mathbb{D}^{p+1}$ guarantees by (5.1) that α cannot be
 330 found again. Repeating this process and considering larger p -periodic MT systems in
 331 each step allows us to find every inverse pole of R .

332 The question of when to stop the above described steps still needs to be considered.
 333 Many popular methods capable of identifying rational functions (for example the
 334 output error model [10]) assume the order of R to be known. If we can assume R has
 335 exactly $p \in \mathbb{N}$ poles, then it is possible to find every inverse pole of R by applying
 336 theorem 2.2 p times. In each step of this process, we can eliminate the inverse pole
 337 α found in the previous iteration by including it in the parameter vector that defines
 338 the current PMT system.

339 One advantage of our proposed pole finding scheme is that it is possible to apply
 340 theorem 2.2 without making any assumptions on the order of R . In this case however,
 341 one has to define a condition on when to stop looking for new inverse poles. We
 342 now propose one such possible stopping condition for the iterative application of
 343 theorem 2.2. Let $\phi_k(z) := \phi_k^{\mathbf{a}_k}$, $\mathbf{a}_k := (a_0, \dots, a_{k-1})$, ($1 \leq k \leq p$) and consider the
 344 p -th Malmquist-Takenaka partial sum of R

$$345 \quad (5.2) \quad S_p R(w) = S_p^{\mathbf{a}} R(w) := \sum_{k=0}^{p-1} c_k \phi_k(w), \quad (R \in \mathcal{A}, w \in \overline{\mathbb{D}}),$$

346 where c_k denote the k -th MTF coefficients.

347 Consider the H_2 norm

$$348 \quad \|f\|_{H_2} = \sqrt{\langle f, f \rangle} \quad (f \in H_2(\mathbb{D}),$$

349 induced by the H_2 scalar product defined in (2.7). Clearly, if $\mathbf{a}_p := (a_0, \dots, a_{p-1})$
 350 exactly matches the inverse poles of R , then $\|R - S_p R\|_{H^2} = 0$ is also true, therefore
 351 we can stop the iteration once the H^2 norm of $R - S_p R$ is zero. Since the inverse
 352 poles of R can be simple and are contained in $A = \{\alpha_0, \dots, \alpha_{p-1}\} \subset \mathbb{D}^p$, the rational
 353 function R belongs to the subspace spanned by $\phi_0, \dots, \phi_{p-1}$. Thus, $\|R - S_p R\|_{H^2} = 0$
 354 indicates, that for the parameter vector generating the partial sum $S_p R$, we have
 355 $\mathbf{a}_p = (\alpha_0, \dots, \alpha_{p-1})$.

356 The steps for the k -th iteration of the proposed pole finding scheme can be sum-
 357 marized as follows.

- 358 1. Identify $\alpha_{k-1} \in A$, by applying theorem 2.2. Let the PMT system involved in
 359 the application of the theorem be generated by $\mathbf{a}_k = (\alpha_0, \dots, \alpha_{k-2}, a) \subset \mathbb{D}^k$,
 360 where $\alpha_0, \alpha_1, \dots, \alpha_{k-2} \in A$.
- 361 2. Use the newly identified α_{k-1} inverse pole to construct the parameter vec-
 362 tor $\mathbf{b}_k := (\alpha_0, \alpha_k, \dots, \alpha_{k-1}) \in \mathbb{D}^k$. Construct the PMT system $\phi_j(z) :=$
 363 $\phi_j^{\mathbf{b}_k}(z)$ ($j = 0, \dots, k-1$).
- 364 3. Consider the orthogonal projection of R onto the subspace spanned by
 365 $\phi_0^{\mathbf{b}_k}, \phi_1^{\mathbf{b}_k}, \dots, \phi_{k-1}^{\mathbf{b}_k}$. This projection can be expressed by the formula in (5.2).
 366 The error of the projection is given by $\|R - S_k R\|_{H^2}$. If this error is zero,
 367 then we have successfully found every inverse pole of R (hence R is completely
 368 contained in the subspace), otherwise increase k and repeat the above steps.

369 In practice, we have to consider a discrete version of the problem. That is, suppose
 370 that instead of R , we only have access to the vector $\mathbf{r} \in \mathbb{C}^N$ ($N \in \mathbb{N}$), where the
 371 components of \mathbf{r} are discrete samplings of R on \mathbb{T} . We may use an equidistant sampling
 372 of \mathbb{T} , or the discrete point set defined in (4.6). If we consider an equidistant sampling,
 373 we have to approximate the integrals $\langle \phi_{\nu_k}^{\mathbf{a}}, R \rangle$ using a numerical quadrature such
 374 as the trapezoid rule when applying theorem 2.2. This approach is quick, however
 375 it introduces numerical errors especially for small N . Instead of this approach, we
 376 can also use the discrete scalar product and discrete orthogonal PMT systems as
 377 discussed in section 4. These allow us more precise computations from a numerical
 378 point of view. In this case however, we have to consider that each application of
 379 theorem 2.2 requires the calculation of the sampling points (4.6) as we modify the
 380 parameter vector defining the PMT system in each iteration of the proposed method.
 381 Thus, using discrete orthogonal PMT systems can increase computational cost. We
 382 note that since the error $\|R - S_p R\|_{H_2}$ depends heavily on R , many signal processing
 383 applications [8, 19] use normalized variations of it. In this work, we propose the use
 384 of percent root mean squared difference (PRD) (see e.g. [19]) to describe the error of
 385 the projection $S_p R$:

$$386 \quad (5.3) \quad PRD(\mathbf{a}) := \sqrt{\frac{\|R - S_p^{\mathbf{a}} R\|_{H_2}^2}{\|R\|_{H_2}^2}} \cdot 100.$$

387 The use of the PRD score allows us to express the error of the approximation with
 388 percentages, thus we can construct a stopping condition for the proposed method
 389 that is usable for any R . In our future work we also plan to explore alternative stop-
 390 ping criteria suited for specific applications, however our experiments (see section 6)

391 demonstrate the usefulness of the proposed approach (5.3). In a computer imple-
 392 mentation of the proposed method, the norm $\|\cdot\|_{H_2}$ is replaced by the $\|\cdot\|_2$ vector
 393 norm, if R was sampled in an equidistant fashion or the norm induced by (4.9), if R
 394 was sampled on (4.6).

395 In algorithm 5.1 we summarize the steps of the proposed inverse pole identifica-
 396 tion and elimination approach. Algorithm 5.1 should not be considered a pseudo-code,
 397 rather a summary of the different steps needed to find the inverse poles of R . In this
 398 formulation we assumed the order of R to be unknown and relied on the above de-
 399 scribed exit condition to stop the iteration. For a more thorough consultation on the
 400 implementation, we refer to our MATLAB implementation of the proposed method [9].

Algorithm 5.1 Generalized Bernoulli's method to find every inverse pole

Obtain \mathbf{r} , a sampling of R on \mathbb{T} .

Let $PRD = 100$ and the exit condition $\varepsilon \in [0, 100]$.

Let $p = 1$.

Let \mathbf{a} and \mathbf{b} be empty vectors.

while $\varepsilon < PRD$ **do**

If $p = 1$, then let $\mathbf{a} = (a) \in \mathbb{D}$. If $p > 1$, then let $\mathbf{a} = (\mathbf{b}_1, \dots, \mathbf{b}_{p-1}, a) \in \mathbb{D}^p$,
 where \mathbf{b}_k denotes the k -th component of the vector \mathbf{b} . In either of these cases, a
 strategy to choose a is given in section 6.

Obtain α_{p-1} by applying theorem 2.2 with \mathbf{a} . The practical application of theo-
 rem 2.2 is discussed in section 6.

If $p = 1$, then let $\mathbf{b} = (\alpha_{p-1}) \in \mathbb{D}$. If $p > 1$, then let $\mathbf{b} = (\mathbf{b}_1, \dots, \mathbf{b}_{p-1}, \alpha_{p-1}) \in \mathbb{D}^p$.
 The vector \mathbf{b} contains the already found inverse poles.

Calculate a discrete version of the projection $S_p^{\mathbf{b}}R$.

Let $PRD = PRD(\mathbf{b})$, where the function $PRD(\mathbf{b})$ is defined in (5.3).

Let $p = p + 1$.

end while

401 **6. Numerical considerations.** In this section we consider some practical prob-
 402 lems that arise when we implement the proposed pole finding scheme. Namely, we
 403 investigate the behavior of the ratios in (2.15), when we can only calculate the MTF
 404 coefficients up to some finite index. In addition, we propose a strategy to choose the
 405 parameter vector $\mathbf{a} \in \mathbb{D}^p$ that defines the Malmquist-Takenaka system in the p -th step
 406 of algorithm 5.1.

407 For a function $f \in H^2(\mathbb{D})$, the modulus of the periodic Malmquist-Takenaka
 408 Fourier coefficients $c_n^{\mathbf{a}} := \langle \phi_n^{\mathbf{a}}, f \rangle$ tends quickly to zero if $n \rightarrow \infty$. This behavior
 409 means that as k increases, calculating the ratios

$$410 \quad (6.1) \quad s_k(\mathbf{a}) := c_{\nu_k}^{\mathbf{a}} / c_{\nu_k - 1}^{\mathbf{a}}$$

411 incurs large numerical errors. On the other hand, considering ratios where the indices
 412 ν_k are too small, the values (6.1) may not approximate the limit well. This problem
 413 is illustrated in Figure 5.

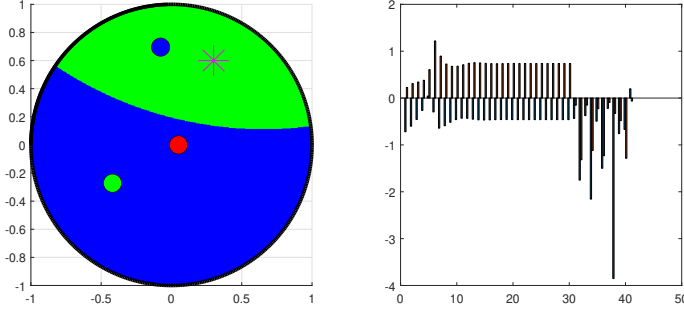


Fig. 5: LEFT: Inverse poles of R (small circles) and the parameter of the PMT system (star) for $p = 1$. RIGHT: Real and imaginary parts of the ratios (6.1). If the index k is too small, the ratios oscillate, and if it is too large numerical errors begin to appear.

414 In order to select the ratios which approximate the limit $s(\mathbf{a}) := \lim_{k \rightarrow \infty} s_k(\mathbf{a})$
 415 closely enough, we have to find an interval $J = [k, k + \ell]$ of indices, where $s_k(\mathbf{a})$ exhibit
 416 "near constant" behavior. To do this, we propose to measure the oscillation in the
 417 window J by

$$418 \quad (6.2) \quad \omega(J, \mathbf{a}) := \max_{i, j \in J} |s_i(\mathbf{a}) - s_j(\mathbf{a})|.$$

419 For a fixed \mathbf{a} we can approximate the limit of (6.1) using

$$420 \quad (6.3) \quad \omega^*(\mathbf{a}) = \min_J \omega(J, \mathbf{a}) = \omega([m^*, m^* + \ell], \mathbf{a}), \quad s(\mathbf{a}) \approx s_{m^*}(\mathbf{a}).$$

421 As shown in section 2, the inverse pole can be recovered by

$$422 \quad \alpha_0 = Q^{\mathbf{a}}(s(\mathbf{a})),$$

423 where the mapping $Q^{\mathbf{a}}$ is defined in (2.17). Consequently, the error formula (2.20)
 424 given as

$$425 \quad |Q^{\mathbf{a}}(s(\mathbf{a})) - Q^{\mathbf{a}}(s_{m^*}(\mathbf{a}))| \leq M(\mathbf{a}) \cdot |s(\mathbf{a}) - s_{m^*}(\mathbf{a})|$$

426 can be used to estimate the error of the reconstruction, where $M(\mathbf{a})$ is defined in (2.21).
 427 Unfortunately, in practice we cannot calculate the exact value of $s(\mathbf{a})$. In this work,
 428 we approximate the error $|s_k(\mathbf{a}) - s(\mathbf{a})|$ with $|s_k(\mathbf{a}) - s_{k+l}(\mathbf{a})|$ ($l, k \in \mathbb{N}$). We now
 429 proceed to show that this error decreases quickly and is therefore appropriate for
 430 most practical cases. By theorem 2.2, there exists $0 \leq q < 1$ and $M \in \mathbb{R}$ for which

$$431 \quad (6.4) \quad |s_k(\mathbf{a}) - s(\mathbf{a})| \leq M \cdot q^k \quad (k \in \mathbb{N}).$$

432 From this, we get

$$433 \quad (6.5) \quad \begin{aligned} |s_k(\mathbf{a}) - s_{k+l}(\mathbf{a})| &\leq M(q+1) \cdot (q^k + q^{k+1} + \dots) \\ &= M \cdot q^k \cdot \frac{1+q}{1-q} \quad (l > 0). \end{aligned}$$

434 By (6.5), the proposed practical error estimate $|s_k(\mathbf{a}) - s_{k+l}(\mathbf{a})|$ has the same order
 435 of decay as $|s_k(\mathbf{a}) - s(\mathbf{a})|$ and can be used in applications. If q is close to 1, then the
 436 proposed estimate is not as reliable, however we did not see a large difference between
 437 the proposed estimate $|s_k(\mathbf{a}) - s_{k+l}(\mathbf{a})|$ and the actual error (6.4) in our experiments.
 438 Finally, we remark, that the MTF coefficients contain information about *every pole*
 439 of R . Methods for the $p = 1$ case have already been developed, where multiple
 440 poles are identified using the PMT expansion of R with $p = 1$ [13]. Based on our
 441 numerical experiments, we conjecture that if we approximate $s(\mathbf{a})$ with $s_n(\mathbf{a})$, where
 442 $s_n(\mathbf{a})$ falls into a "relatively constant" part of the sequence $s_k(\mathbf{a})$, then $Q^{\mathbf{a}}(s_n(\mathbf{a}))$ will
 443 approximate *one of the inverse poles* of R (not necessarily α_0). We plan to study this
 444 phenomenon and formalize our findings in a future work.

445 We found in our experiments that estimating the value of the error formula (2.20)
 446 with

$$447 \quad (6.6) \quad E(m^*, \mathbf{a}) := M(\mathbf{a}) \cdot \omega^*(\mathbf{a})$$

448 suffices whenever the order of R is not too large.

449 Next, we propose an approach to automatically choose the parameter $\mathbf{a} \in \mathbb{D}^p$
 450 in the p -th step of algorithm 5.1. By the error formulas (2.20) and (6.6) it is clear
 451 that the error of the inverse pole reconstruction depends heavily on the parameter
 452 vector \mathbf{a} . A poor choice of \mathbf{a} can mean that the sequence (6.1) converges slowly. This
 453 phenomenon is illustrated in figure 6. In this sense we can find a good parameter
 454 vector \mathbf{a} by minimizing the function

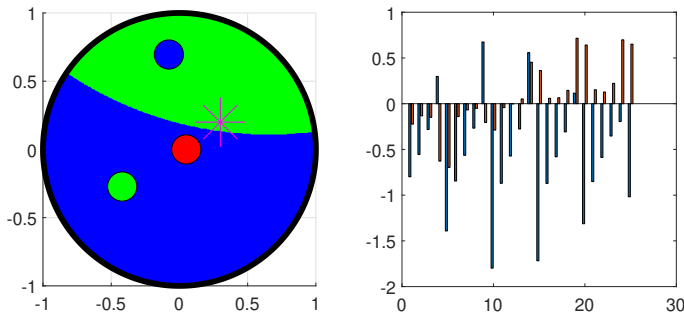


Fig. 6: LEFT: Inverse poles of R (small circles) and the parameter of the PMT system (star) for $p = 1$. RIGHT: Real and imaginary parts of the ratios (6.1). If the parameters of the PMT system lie close to the border of the Voronoi cells discussed in section 3, then convergence of the ratios (6.1) is slow.

$$455 \quad (6.7) \quad E_{m^*}(\mathbf{a}) := E(m^*, \mathbf{a}),$$

456 where, for any given \mathbf{a} , the index m^* is determined by (6.3). Minimizing (6.7) leads
 457 to a nonlinear optimization problem. We note that in the p -th step of the algorithm,
 458 the first $p - 1$ components of \mathbf{a} are fixed (they are the inverse poles reconstructed in
 459 previous steps), therefore we only have to find a single $a_p \in \mathbb{D}$ parameter which mini-
 460 mizes (6.7). In this work we considered two algorithms to solve the above mentioned

461 optimization problem. In the first case, we considered 10 random $a \in \mathbb{D}$ at each step
 462 of the algorithm and selected the one for which (6.7) was minimal. In the second case,
 463 we used the hyperbolic variant of the Nelder-Mead simplex method [23]. The Nelder-
 464 Mead method [26] can be used to solve nonlinear optimization problems. It applies
 465 successive geometric transformations to a simplex, whose vertices represent the cur-
 466 rent state of the minimization. The applied transformations depend on the objective
 467 function's values at the vertices. The hyperbolic Nelder-Mead algorithm introduced
 468 in [23] replaces these geometric transformations with their hyperbolic variants. When
 469 minimizing the objective (6.7) this is useful, as it naturally ensures all components of
 470 the vector \mathbf{a} remain strictly inside \mathbb{D} . We note that the proof of convergence can only
 471 be given in simple cases, even in the original variant of the Nelder-Mead method (see
 472 e.g. [23]). Despite this, it remains a popular minimization method based on empirical
 473 evidence and the results of our experiments also confirm its usefulness for the problem
 474 stated above. In particular, our below numerical results demonstrate the effectiveness
 475 of the proposed method when used with the above mentioned optimization schemes.

476 We created a MatLab implementation of the proposed methods which can be
 477 accessed at [9]. To calculate periodic MT systems and the corresponding coefficients,
 478 we relied on the library introduced in [21]. Below, we provide an example to demon-
 479 strate the effectiveness of the proposed algorithm. We consider the rational function
 480 R given by the inverse poles $A := \{0.3 + 0.4i, -0.5 - 0.4i, 0.7 - 0.3i\}$ and zeros
 481 $z_0 = 0.8 + 0.4i, z_1 = 0.8 + 0.4i$. We choose parameter \mathbf{a} in each step of the algorithm
 482 by minimizing (6.7) by the above described optimization methods. The results for
 483 this example can be found in table 1. The rows of the table represent the iterations
 484 of algorithm 5.1.

	Hyperbolic Nelder Mead		Monte Carlo optimization	
Step	$ \alpha_0 - Q^{\mathbf{a}}(s_{m^*}(\mathbf{a})) $	$E_{m^*}(\mathbf{a})$	$ \alpha_0 - Q^{\mathbf{a}}(s_{m^*}(\mathbf{a})) $	$E_{m^*}(\mathbf{a})$
1	$5.03 \cdot 10^{-12}$	$4.26 \cdot 10^{-11}$	$1.86 \cdot 10^{-9}$	$1.83 \cdot 10^{-8}$
2	$5.02 \cdot 10^{-15}$	$3.73 \cdot 10^{-14}$	$5.39 \cdot 10^{-11}$	$2.71 \cdot 10^{-10}$
3	$1.48 \cdot 10^{-15}$	$1.35 \cdot 10^{-14}$	$1.04 \cdot 10^{-15}$	$1.28 \cdot 10^{-14}$

Table 1: Results for the above described example problem. The columns $|\alpha_0 - Q^{\mathbf{a}}(s_{m^*}(\mathbf{a}))|$ and $E_{m^*}(\mathbf{a})$ refer to the actual error of the reconstruction and the error estimate with the optimized parameters (6.6).

485 In table 2, we present the results of a larger simulation. In this case, we construc-
 486 ted rational functions of the form $R(z) := \sum_{k=1}^M c_k \cdot r_{\alpha_k}(z)$ ($c_k \in \mathbb{C}, \alpha_k \in \mathbb{D}, z \in \overline{\mathbb{D}}$),
 487 where the coefficients c_k and the inverse poles α_k were chosen randomly. We con-
 488 ducted 100 such experiments for each $M = 1, \dots, 5$, with table 2 showing the mean
 489 distance of the estimates from the actual inverse poles of R . The proposed method
 490 was applied with hyperbolic Nelder-Mead optimization. To simplify the evaluation of
 491 the results, the number of poles of R was assumed to be known in these experiments.
 492 In table 2, the error values for each M are given by

$$493 \quad (6.8) \quad \text{Err}(M) := \frac{1}{100} \sum_{k=1}^{100} \left(\frac{1}{M} \sum_{j=0}^{M-1} |\alpha_{k,j} - Q_{k,j}^{\mathbf{a}}(s_{m^*}(\mathbf{a}))| \right),$$

494 where $\alpha_{k,j}$ denotes the j -th inverse pole of the k -th rational function which is defined

495 by M poles and $Q_{k,j}^a(s_{m^*}(\mathbf{a}))$ denotes the estimate of $\alpha_{k,j}$ produced by our proposed
 496 method. The results in table 2 show, that our proposed method can be used reliably
 497 to find the inverse poles of R . For elementary rational functions (when $M = 1$),
 498 the reconstruction is almost perfect even for a large number of experiments. When
 499 increasing the number of inverse poles that define R , we can see a decrease in precision,
 500 however the average error defined in (6.8) remains in the order of 10^{-5} even if $M = 5$.

501 Finally, we conducted an experiment to measure the effectiveness of the stopping
 502 criteria for our algorithm proposed in section 5. In particular we generated 100
 503 rational functions, each with 5 poles and applied the proposed method to find every
 504 inverse pole. This time however, we did not assume the number of poles to be known
 505 in advance, instead we stopped our iteration once the value of the PRD error (5.3)
 506 became less than $\varepsilon = 10$. We found that the average number of identified inverse
 507 poles throughout the 100 experiments in this case was 4.3. A perfect score could not
 508 be expected, because some inverse poles contribute very little to the energy ($\|R\|_{H_2}$)
 509 of R , however the results show that we can rely on this scheme to accurately identify
 510 most significant inverse poles. We note that lowering the threshold ε increases the
 511 number of identified inverse poles, however it also increases computational cost (as
 512 the algorithm will keep looking for new inverse poles even after the most dominant
 513 ones have been found).

M	1	2	3	4	5
Err(M) see (6.8)	$3.0 \cdot 10^{-16}$	$1.5 \cdot 10^{-6}$	$4.2 \cdot 10^{-7}$	$1.9 \cdot 10^{-5}$	$2.9 \cdot 10^{-5}$

Table 2: Results of a larger experiment with different numbers of poles.

514 In our experiments we found that the proposed algorithm can be used to reliably
 515 identify the inverse poles of rational functions. Even though both described optimiza-
 516 tion methods provided good estimates on the inverse poles, applying the hyperbolic
 517 variant of the Nelder-Mead optimization showed slightly better precision. In the ex-
 518 periment detailed in table 1, the average distance between the estimated and true
 519 inverse poles was $1.68 \cdot 10^{-12}$ with the Nelder-Mead method and $6.39 \cdot 10^{-10}$ if we
 520 used a Monte Carlo approach. Every experiment was conducted using the algorithm
 521 in our implementation [9]. The above results justify using nonlinear optimization al-
 522 gorithms adapted to hyperbolic geometry to minimize (6.7). In our future work, we
 523 plan to experiment using further hyperbolic optimization methods such as [20].

524 **7. Conclusion.** In this work we introduced a generalization of Bernoulli's clas-
 525 sical method of finding the poles of a rational function. The generalization uses
 526 periodic Malmquist-Takenaka Fourier coefficients to construct the sequence of ratios
 527 used by Bernoulli's original algorithm. We generalized the concept of dominant poles
 528 using Blaschke-products and gave a description of the poles which can be found with
 529 the proposed method. Furthermore, we showed that discrete orthogonal Malmquist-
 530 Takenaka systems can also be used with the proposed method. Using our results, we
 531 proposed an iterative algorithm which applies the generalized Bernoulli's method to
 532 find every inverse pole of the rational function. Finally, we proposed a method to
 533 automatically select the parameters of our algorithm by minimizing an intuitive cost
 534 function with different optimization techniques.

535 The proposed method is an interesting generalization of a classical numerical al-
 536 gorithm, which in our opinion is worthy of attention by itself. In addition however, the

537 proposed method exhibits great practical potential in the field of system identifica-
 538 tion. Specifically, in our future work we plan to investigate ways in which to apply the
 539 proposed algorithm to identify the poles of the transfer functions of SISO-LTI (single
 540 input single output, linear time invariant) systems [2]. One promising property of the
 541 proposed algorithm is that the order of the transfer function to be identified need not
 542 be known in advance.

543 Another area of future investigation will be the description of identifiable inverse
 544 poles through convex geometry. As mentioned in section 3, when the generalized al-
 545 gorithm is used with Laguerre (1-periodic Malmquist-Takenaka) Fourier coefficients,
 546 the set of identifiable inverse poles can be given by calculating their so-called para-
 547 cyclic convex hull. This result and further generalizations for the case $p > 1$ will be
 548 considered in our future research.

549 **Acknowledgments.** Project no. C1748701, K146721 have been implemented
 550 with the support provided by the Ministry of Culture and Innovation of Hungary
 551 from the National Research, Development and Innovation Fund, financed under the
 552 NVKDP-2021 and the K.23 "OTKA" funding schemes, respectively. The research was
 553 supported by the European Union within the framework of the National Laboratory
 554 for Autonomous Systems. (RRF-2.3.1-21-2022-00002).

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